

StorPool Storage

#SFD18, @storpool



Introduction to StorPool

Boyan Ivanov, co-founder & CEO #SFD18, @storpool

Best-of-breed block storage software (1)

- 1. Software company scale-out, block storage software
 - a. Primary, flash (SATA/NVMe)
- 2. Not your typical Valley startup
- 3. Doing this before SDS/SDN/SDDC & *MARKETING-DEFINED STORAGE*
- 4. Delivered as a working storage solution on customer's hardware:
 - a. Fully managed: software + 24/7/365 support, SLA, proactive monitoring
 - b. Hardware Compatibility List (HCL) or
 - c. A pre-integrated solution with partners
- 5. SDS 2.0 feature rich shared storage system faster than local SSD-



Best-of-breed block storage software (2)

- 6. Developed from scratch:
 - a. Own on-disk format, protocol, quorum, client, etc, etc.
 - b. Fully distributed, scale-out, online changes of everything, etc.
 - c. Running in production for 6+ years; numerous 1PB+ flash systems; 17 major releases; Global spread of customers
- 7. Target customers companies building public & private clouds:
 - a. Service providers & public clouds
 - b. Enterprises & various private clouds
- 8. Use cases anything block DBs, VM disks, VDI, etc.
- 9. Replacing single-purpose SAN / AFA or other storage software



Best-of-breed block storage software - Example:

NVMe shared storage system with:

- 1. **Latency**: < 100 μs
- Throughput: >1M IOPS per server (scale-out, 10 servers = 10M IOPS @ ~ 300µs)
- Feature rich: API, end-to-end data integrity, self-healing, online everything, thin provisioning, snaps & clones, QoS, backup & DR, etc.
- 4. **Fully managed**: 24/7 support; SLA; proactive monitoring & issue resolution
- 5. **TCO over 3 years**: < \$0.10 /GB provisioned /month; < \$0.002 /IOPS/month

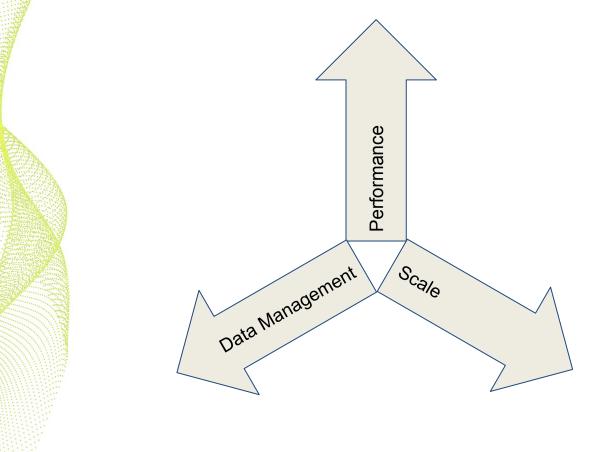




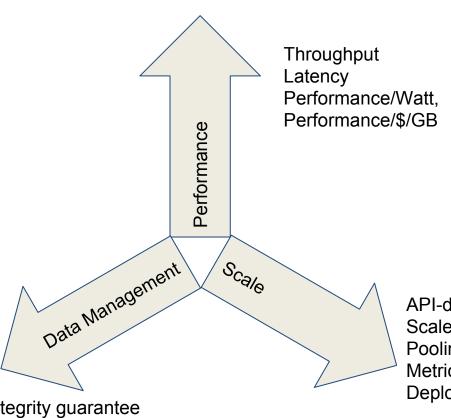


You can't have your cake and eat it or can you?

Boyan Krosnov, Co-founder and CPO #SFD18, @storpool



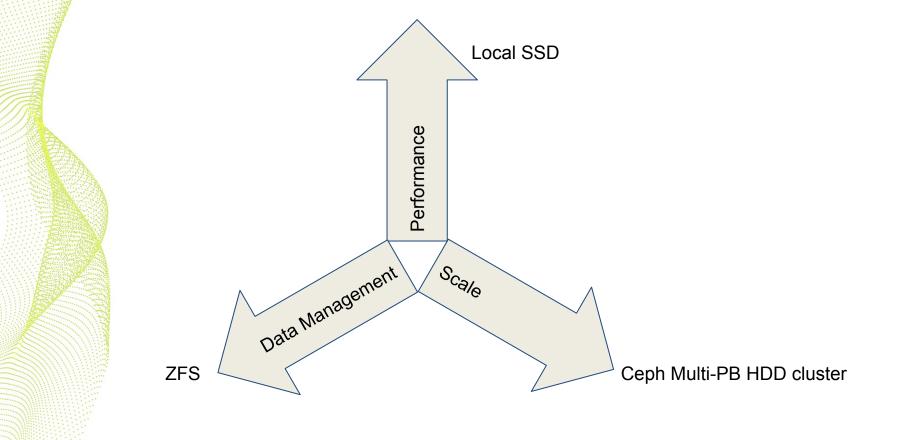




API-driven, integrations Scale by adding nodes/drives Pooling of capacity & performance Metrics collection & Monitoring Deployment automation

End-to-end data integrity guarantee "LUN" per vDisk CoW - Thin provisioning, Snapshots, Clones Multi-site with efficient transport of changes Fast recovery (changed block tracking)







Why performance

Fast storage system = more work done per CPU

The virtualization & cloud promise:

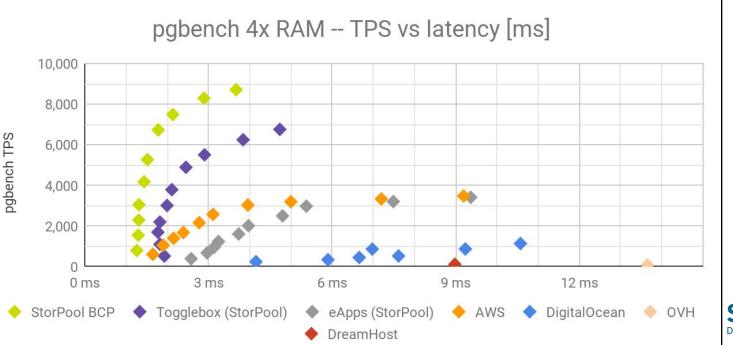
- near-zero overhead
- vDisks as fast as a local SSD are they really?
- efficient consolidation of workloads



Why performance

Fast storage system = more work done per CPU

The virtualization & cloud promise:





Performance with StorPool

>1M IOPS per node >250k IOPS per core (server) >500k IOPS per core (client)

StorPool latency overhead on par with latency of NVMe devices. End-to-end latency approx 2x local NVMe latency.

100k IOPS per NVMe drive with <150 μ s end-to-end latency Writes at QD1 ~ 70 μ s end-to-end





Why Scale

Public and private cloud Mobile and Web apps, SaaS Containers & Microservices DevOps, Infrastructure as code

What is scale:

API-driven, integrations Scale by adding nodes/drives Pooling of capacity & performance Metrics collection & Monitoring Deployment automation



Scale with StorPool

Scale-out architecture

>1PB usable All-SSD & Hybrid clusters in production for years Some customers of StorPool have multiple clusters per location and multiple locations

API control and integrations with Kubernetes, OpenStack, OpenNebula, CloudStack & OnApp

Detailed metrics collection, monitoring.





Why Data Management

Assumed that every storage system has it But many don't

What is data management:

End-to-end data integrity guarantee CoW - Thin provisioning, Snapshots, Clones "LUN" per vDisk Multi-site with efficient transport of changes Fast recovery (changed block tracking)



Data Management with StorPool

End-to-end data integrity protection

4k granularity

- thin provisioning / reclaim
- CoW snapshots, clones
- changed block tracking, incremental recovery and transfer

Multi-site

- connect 2 or more StorPool clusters over public Internet
- send snapshots between clusters for backup and DR
- commonly 100TB backup once per hour



Performance goal: vDisks as fast as Local SSD Scalability goal: Better than Ceph Data Management goal: Better than ZFS

In one system - StorPool



And as conclusion...



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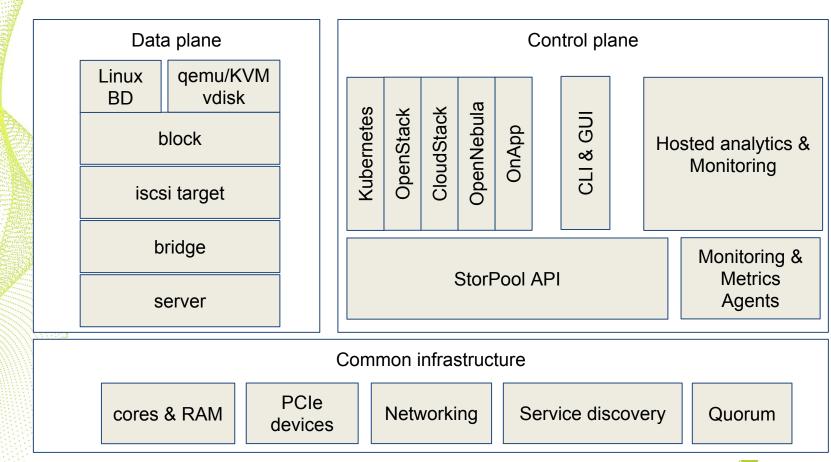






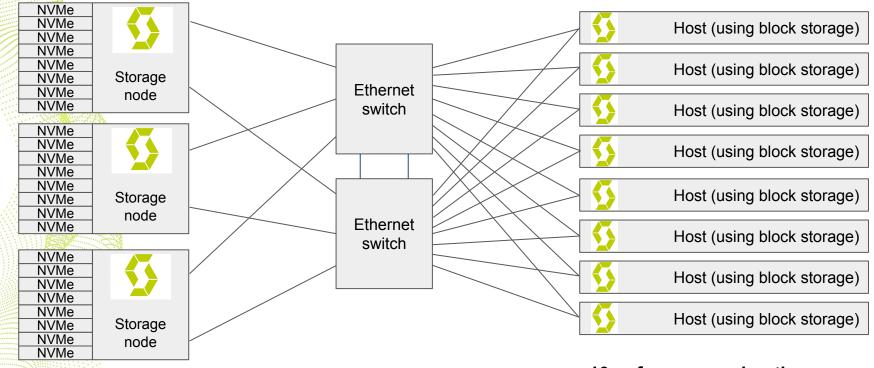
Architecture and Demo

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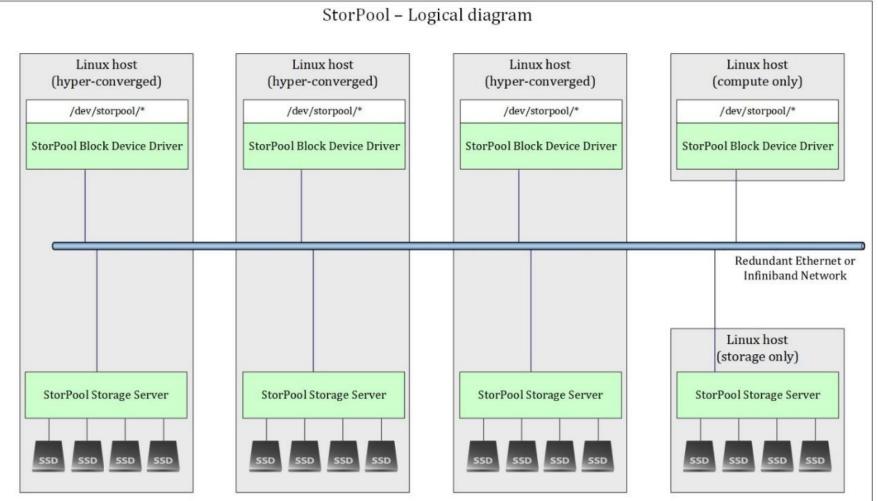


Data plane

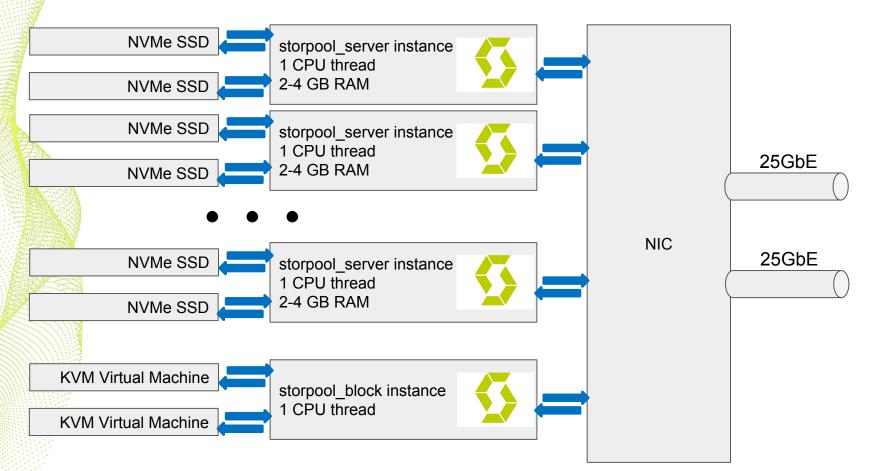


3+ storage nodes Scale-out ... 10s of servers using the storage system





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- Highly scalable and efficient architecture
- Scales up in each storage node & out with multiple nodes



Protection schemes:

- 3 copies on SATA SSDs or NVMe SSDs
- StorPool Hybrids 1+2 or 2+1 lower cost
- 3 copies on HDD
- Erasure coding soon

Writes are sequentialized and coalesced. Under load 1 write above is less than 1 write below.

Writes may go through "journal" write-back devices before "pool". Evolution:

- 3108 RAID controller w/ CacheVault legacy
- fast SSD, including "pool" SSDs current
- Intel Optane drive current
- NVDIMM / PM future



Demo:

- Private cloud use-case
- Basic CLI
 - Analytics
 - UI dashboard



Analytics / metrics collection

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Front-end IOps 14.78K iops		Front-end IO 439 MBps		Front-end latency 388 µs		Back-end IOps 23.1K iops		Back-end IO 637 MBps		Back-end latency 177 μs	
Clients (minut Search	te data)	Clients (seconds data) Search		Servers (minute data) Search		Servers (seconds data) Search		Volumes (minute data) Search		Volumes (seconds da Search	ta)
All client stat		All client stat		All servers stat		All servers stat		Per-volume stat		Per-volume stat	
General client stat		General client stat		General Server Stat		General Server Stat		Top volumes			
Per-host client stat		Per-host client stat		Per-disk backend stat		Per-disk backend stat					
				Per-server backend stat		Per-server backend stat					
CPU (minute Search	data)	CPU (seconds data) Search		Cgroup memory (minute da Search	ata)	Cgroup memory (seconds d Search	lata)	Templates Search		Disks Search	
CPU Stat		CPU Stat		Cgroup memory		Cgroup memory		Template usage		Disk usage	
Per-CPU stat		Per-CPU stat		Cgroup memory per node		Cgroup memory per node		Template usage - internal		Disk usage - internal	
Per-service CPU Stat		Per-service CPU Stat								Single disk usage - internal	
Total CPU Non-SP Stat		Total CPU Non-SP Stat									
Total CPU Stat		Total CPU Stat									

Analytics / metrics collection

🤨 - 🗱 Top volume	✓ Zoom Out > ② Last 6 hours 3										
Number of volumes 5 • Skip volumes 5bf50f01f11731.01393943_context •											
	Top 5 v	olumes by read bytes	Top 5 volumes by reads								
total -	bps	volume	total 🔻	rps	volume						
944.61 GIB	46.96 MBps	5bf7e61075ed01.01242219_root	20.74 MII	960.15 iops	5bf7e61075ed01.01242219_root						
437.41 GIB	21.74 MBps	5c0533bfdee630.68819684_root	10.32 Mil	477.70 iops	shared_hertz_home						
380.50 GIB	18.92 MBps	5c177c4c52de63.97691854_root	8.56 MII	396.57 iops	shared_server28_home						
307.22 GIB	15.27 MBps	5c13ded49ccbe1.46303538_root	8.53 MII	394.84 iops	shared_server25_home						
265.24 GIB	13.18 MBps	shared_server21_home	8.37 Mil	387.51 iops	5c0533bfdee630.68819684_root						
	Top 5 vo	lumes by written bytes	Top 5 volumes by writes								
total -	bps	volume	total +	wps	volume						
187.87 GIB	9.35 MBps	shared_server29_home	2.94 MII	135.99 lops	5c46e605e57d55.22476738_ssd						
81.31 GIB	4.04 MBps	shared_server24_home	2.38 MII	110.12 lops	shared_hertz_var						
68.48 GIB	3.40 MBps	shared_server28_localbkp	1.83 MII	84.52 lops	shared_server29_home						
64.48 GIB	3.21 MBps	mailvps_home	1.53 Mil	70.63 iops	shared_hopkins_var						
63.39 GIB	3.15 MBps	shared_server23_localbkp	1.29 Mil	59.76 iops	shared_dio_var						

+ ADD ROW

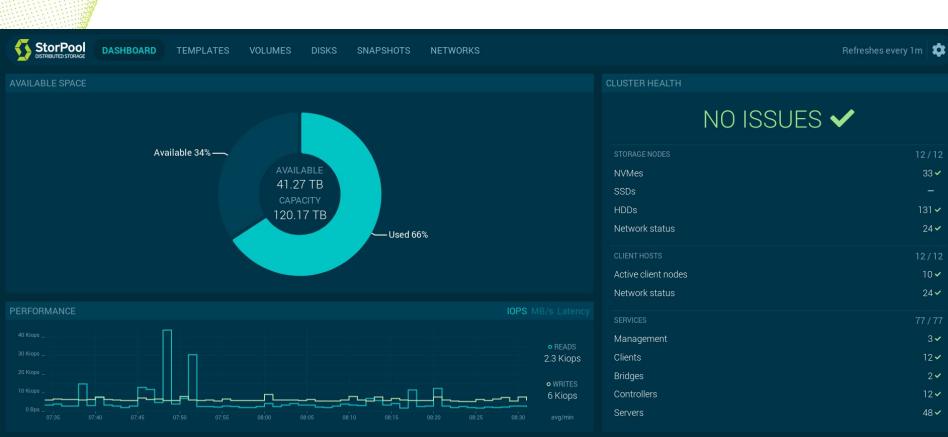
Analytics / metrics collection

🔄 🗸 📗 Per-volume stat 🗸 🖻 🌼

IOps for shared_hertz_home Bps for shared_hertz_home 100 MBps 80 MBps 60 MBps 40 MBps 20 MBps 0 Bps - Writes \$tag_volume Avg: 52 iops - Reads \$tag_volume Avg: 477 iops - Written shared_hertz_home_Avg: 774 kBps - Read shared_hertz_home_Avg: 11.0 MBps Average wait time for shared_hertz_home Queue depth for shared_hertz_home 0.05 - Write shared hertz home Avg: 304 µs - Read shared hertz home Avg: 238 µs - mean shared hertz home Avg: 0



✓ Zoom Out > ④ Last 15 minutes Refresh every 10s





131 -

24~

48 -





Performance demo

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Intel® Data Center Builders

Intel® Builders Construction Zone



"The new HCI industry record: 13.7 million IOPS with Windows Server 2019 and Intel® Optane[™] DC persistent memory"

https://blogs.technet.microsoft.com/filecab/2018/10/30/windows-serve r-2019-and-intel-optane-dc-persistent-memory/



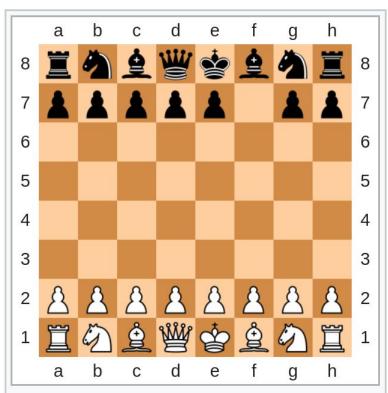
Microsoft's HCI setup

- 12 nodes, each with:
 - 384 GiB (12 x 32 GiB) DDR4 2666 memory
 - 2 x Intel® Xeon® Scalable processor
 - 1.5 TB Intel® Optane[™] DC persistent memory as cache
 - 32 TB NVMe (4 x 8 TB Intel® DC P4510) as capacity
 - 2 x Mellanox ConnectX-4 25 Gbps w/RDMA
 - For the best performance, every VM runs on the server node that owns the volume where its VHDX file is stored.
- S2D, Hyper-V, Windows Server 2019



Handicaps (or "odds") in chess are variant ways to enable a weaker player to have a chance of winning against a stronger one. There are a variety of such handicaps, such as material odds (the stronger player surrenders a certain piece or pieces), extra moves (the weaker player has an agreed number of moves at the beginning of the game), extra time on the chess clock, and special conditions (such as requiring the odds-giver to deliver checkmate with a specified piece or pawn). Various permutations of these, such as "pawn and two moves", are also possible.

Handicaps were quite popular in the 18th and



Initial setup for *pawn and move*: Black starts without the f-pawn; the weaker player (White) moves first

Microsoft's HCI setup

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- StorPool, KVM, CentOS 7



- 12 nodes, each with:
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 - 32 TB NVMe (4 x 8 TB Intel® DC P4510) as capacity
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- StorPool, KVM, CentOS 7



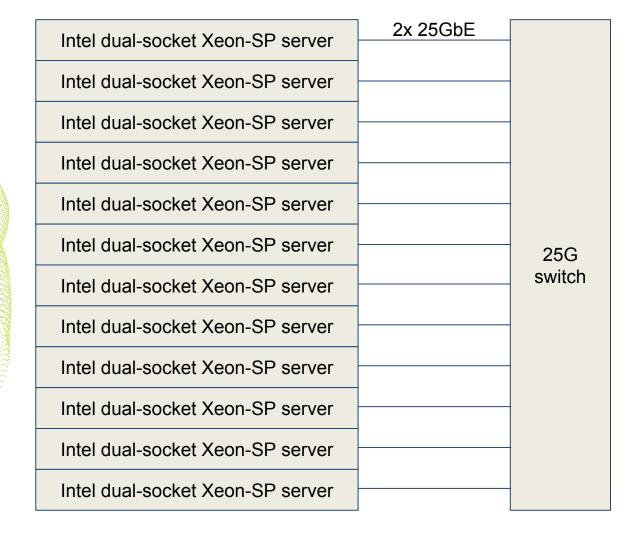
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 - ADD: Intel XXV710-DA2 dual-port 25 Gbps
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- StorPool, KVM, CentOS 7



StorPool's HCI setup

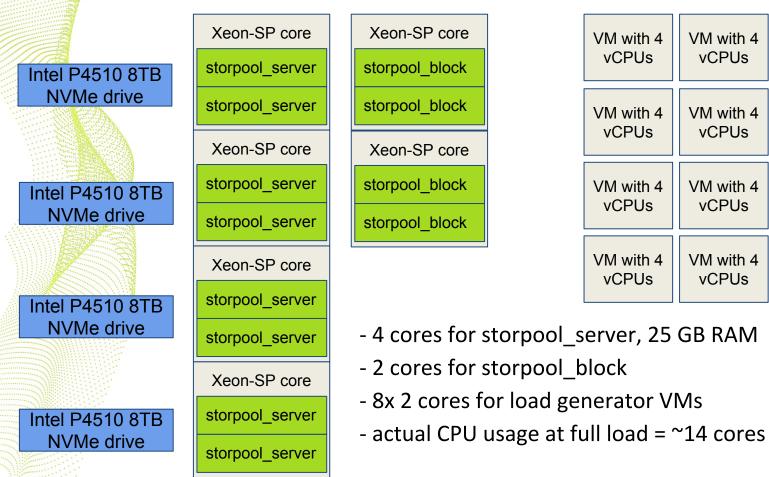
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 - 2 x Mellanox ConnectX-4 25 Gbps w/RDMA REMOVE
 - ADD: Intel XXV710-DA2 dual-port 25 Gbps
 - For the best performance, every VM runs on the server node that owns the volume where its VHDX file is stored. No! 100% remote.
- StorPool, KVM, CentOS 7







Resource usage in each node





Summary of results

result		parameters	comment
13.8 M IOPS	Random read	4k qd 96x64	1.15M IOPS per node
5.5 M IOPS	Random R/W	70/30 4k qd 96x40	183k read/writes /s per drive
2.5 M IOPS	Random write	4k qd 96x40	156k writes /s per drive
64.6 GB/s	Sequential read	bs 128k qd 96x16	
20.8 GB/s	Sequential write	bs 128k qd 96x16	
70 µs	Write latency	bs 4k qd 1	



Active set discussion

Active set in Microsoft / Hyper-V / S2D 312 VMs * 10 GiB each = 2.9 TB (15% of 19.8 TB Optane DC persistent memory)

Active set in StorPool / KVM 96 VMs * 500 GiB each = 44.7 TB (38% of 116 TB system capacity on P4510 NVMes)

storpool_server memory 12 servers * 8 instances * 3.1 GiB = 277 GB







Case studies

Boyan Ivanov, co-founder & CEO #SFD18, @storpool

Case study 1: NVMe-powered VDI cloud - requirement

The need:

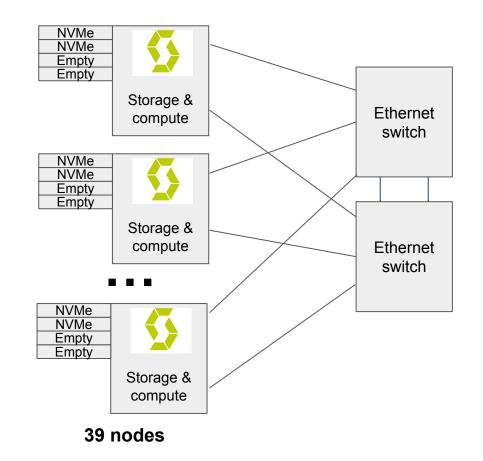
- Fast & cost efficient VDI as a service
- Minimal CPU footprint of software layers
 - Latency: as low as possible

The solution:

- First stage: 39 servers, running Hyper-converged (Compute+storage)
- KVM + StorPool + CloudStack
- 2 CPU cores for StorPool (server & client)
- Just 2 NVMe per server, ~90 TB usable
- 2 x 25 GbE Ethernet



Case study 1: NVMe-powered VDI cloud - diagram







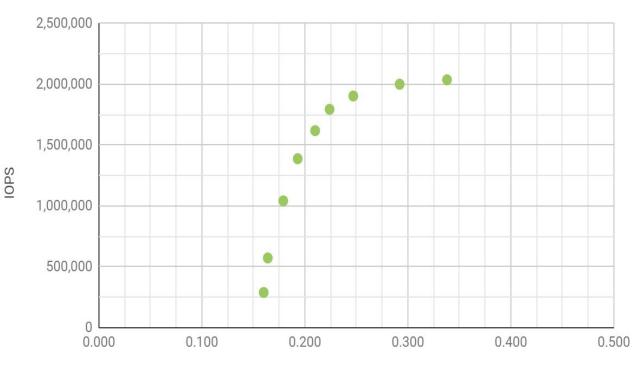
Case study 1: NVMe-powered VDI cloud - metrics

Test	Block size	Queue Depth	Result	Unit
Latency read	4k	1	0.147	ms
Latency write	4k	1	0.111	ms
Random read	4k	64	6,829,218	IOPS
Random read/write	4k	64	1,904,408	
Random write	4k	64	980,619	IOPS
Sequential read	1 M	64	72,072	MB/s
Sequential write	1M	64	20,716	



Case study 1: NVMe-powered VDI cloud - IOPS vs. latency

IOPS vs. ms



Case study 2: Imperia Online - requirement

The need:

- Massively multiplayer online real-time strategy game (MMORTS)
- 40 million users

Game responsiveness & latency is most important

Zero downtime target

The solution:

- KVM + StorPool + OpenNebula
- Multiple StorPool clusters
 - Each cluster: 5 storage nodes & 40 hypervisors
- Each storage node: 4x SATA SSDs
- 2 x 10 GbE Ethernet



Case study 2: Imperia Online



Username

Password

Login Lost password

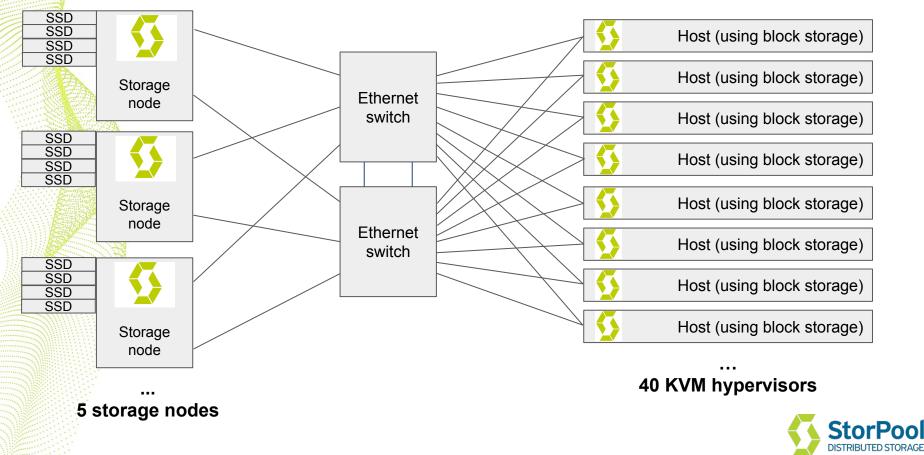
G Google

HALL OF FAME

VIEW

Register Now

Case study 2: Imperia Online - diagram



Case study 2: Imperia Online - Metrics

- 100% uptime
 - Constant < 1ms real life latency
 - Page loading time was reduced from 200-300 milliseconds per page to 75-100 milliseconds.
 - This made Imperia Online 4 times faster than its biggest competitor.

Financial Times study on page load time (28 days measurement):

- 1 second slower: -4.57% in USD revenue
- 3 seconds slower: -7.89% in USD revenue





Thank you!

Boyan I. / Boyan K.

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